

# **Network Alignment, Graph Similarity, and Sports Ranking**

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*Southeastern Ranking Workshop  
College of Charleston*

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In collaboration with  
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Funded by Microsoft Live Labs Fellowship and the Library of Congress.

# Administrivia

“Ladies and gentleman, please take your seat and fasten your seat belt. And also make sure your seat back and folding trays are in their full upright and locked position.”

“At this time, we request that all cellular phones, pagers, radios and remote controlled toys be turned off for the full duration of the flight, as these items might interfere with the navigational and communication equipment on this aircraft.”

PAPER [arxiv.org/abs/0907.3338](https://arxiv.org/abs/0907.3338)

Google “network alignment gleich”

CODE [stanford.edu/~dgleich/publications/2009/netalign](https://stanford.edu/~dgleich/publications/2009/netalign)

Google “netalign gleich”

## Desserts

**URI:** <<http://id.loc.gov/authorities/sh85037243#concept>>

**Type:** Topical Term

**Broader Terms:**

- Confectionery

**Narrower Terms:**

- Ambient desserts
- Banana splits
- Charlottes (Desserts)
- Chocolate desserts
- Frozen desserts
- Ice cream cones
- Mousses
- Puddings
- Refrigerated desserts
- Sundaes
- Whipped toppings

**LC Classification:** TX773

**Created:** 1986-02-11

**Last Modified:** 1988-01-15 17:36:44

# Category:Desserts

From Wikipedia, the free encyclopedia

*The main article for this **category** is **[dessert](#)**.*

**Desserts** are **sweet foods** eaten purely for pleasure, typically at the end of a **meal**.



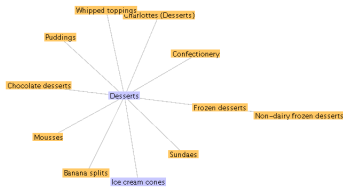
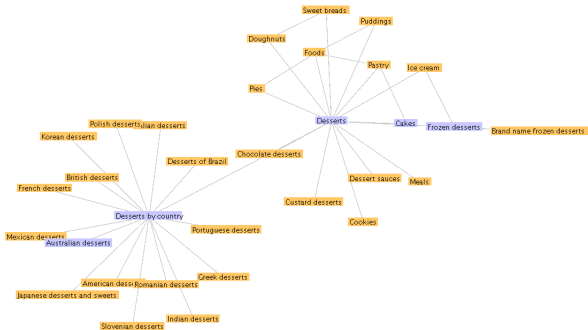
Wikimedia Commons has media related to: ***Desserts***

## Subcategories

This category has the following 16 subcategories, out of 16 total.

- |  |  |   |
|--|--|---|
| <b>*</b>                                       | <b>C cont.</b>                             | <b>P</b>                                |
| ▪ <a href="#">[+] Desserts by country</a> (20) | ▪ <a href="#">[+] Cookies</a> (1)          | ▪ <a href="#">[+] Pastry</a> (4)        |
| <b>B</b>                                       | ▪ <a href="#">[+] Custard desserts</a> (0) | ▪ <a href="#">[+] Pies</a> (9)          |
| ▪ <a href="#">[+] Brand name desserts</a> (3)  | <b>D</b>                                   | ▪ <a href="#">[+] Puddings</a> (2)      |
| <b>C</b>                                       | ▪ <a href="#">[+] Dessert sauces</a> (0)   | <b>S</b>                                |
| ▪ <a href="#">[+] Cakes</a> (0)                | ▪ <a href="#">[+] Doughnuts</a> (1)        | ▪ <a href="#">[+] Sweet breads</a> (2)  |
| ▪ <a href="#">[+] Chocolate desserts</a> (0)   | <b>F</b>                                   | <b>μ</b>                                |
| ▪ <a href="#">[+] Confectionery</a> (9)        | ▪ <a href="#">[+] Frozen desserts</a> (2)  | ▪ <a href="#">[+] Dessert stubs</a> (1) |
|  | <b>I</b>                                   |   |
|  | ▪ <a href="#">[+] Ice cream</a> (5)        |   |

# Library of Congress vs. Wikipedia



**Are they similar?**

*Note Only a bit contrived, but I didn't have to work too hard to find the example*

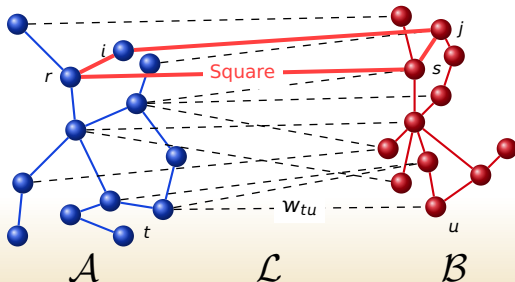
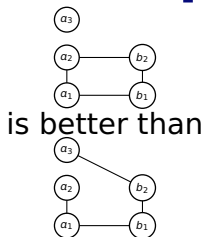
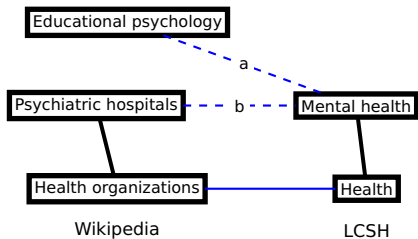
# **Network alignment**

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Network alignment

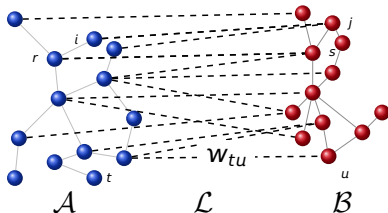
Graph similarity

# Alignment and overlap



**Maximize squares/overlap in 1-1 matching**

# LCSH/Wikipedia alignment



A	LCSH	297,266 vertices, 248,230 edges
B	Wikipedia	205,948 vertices, 382,353 edges
L	links	4,971,629 edges

## MAX-WEIGHT BIPARTITE MATCHING

Overlap 2,346

Weight 60,120 (106,294)

**Using structure should do better!**

Note Links generated by a text search Lucene and weighted with Soft TF/IDF scoring



# Max-weight matching math

## Variables, Data

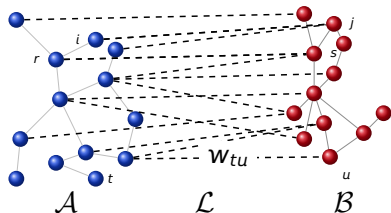
$x_i$  = edge indicator

$w_i$  = weight of edges

$$e_i \in \mathcal{L}$$

$$e_i = (t, u)$$

$$w_i = w_{tu}$$



## Problem

$$\text{maximize}_{x_i} \sum_{i: e_i \in \mathcal{L}} x_i w_i$$

$$\text{subject to} \sum_{i: e_i = (t, u) \in \mathcal{L}} x_i \leq 1, \quad \forall t \in \mathcal{A}$$

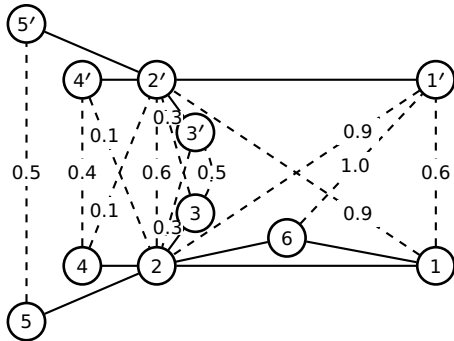
$$\sum_{i: e_i = (t, u) \in \mathcal{L}} x_i \leq 1, \quad \forall u \in \mathcal{B}$$

$$x_i \in \{0, 1\}$$



$$\begin{aligned} &\text{maximize}_{\mathbf{x}} \quad \mathbf{w}^T \mathbf{x} \\ &\text{subject to} \quad \mathbf{A} \mathbf{x} \leq \mathbf{e} \\ &\quad \quad \quad x_i \in \{0, 1\} \end{aligned}$$

# Example



- $e_1=(2,2')$
- $e_2=(2,1')$
- $e_3=(2,3')$
- $e_4=(2,4')$
- $e_5=(1,2')$
- $e_6=(1,1')$
- $e_7=(3,2')$
- $e_8=(3,3')$
- $e_9=(4,2')$
- $e_{10}=(4,4')$
- $e_{11}=(5,5')$
- $e_{12}=(6,1')$

$$\mathbf{W} = \begin{pmatrix} 0.6 \\ 0.9 \\ 0.3 \\ 0.1 \\ 0.9 \\ 0.6 \\ 0.3 \\ 0.5 \\ 0.1 \\ 0.4 \\ 0.5 \\ 1.0 \end{pmatrix}$$

$$\mathbf{A} = \begin{pmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

# Matching and overlap

Squares produce overlap  $\rightarrow$  bonus for some  $x_i$  and  $x_j \rightarrow \sum x_i x_j$

## Variables, Data

$x_i$  = edge indicator

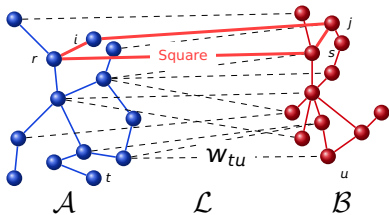
$e_i \in \mathcal{L}$

$w_i$  = weight of edges

$e_i = (t, u)$

$S_{ij}$  squares in  $\mathcal{S}$

$w_i = w_{tu}$

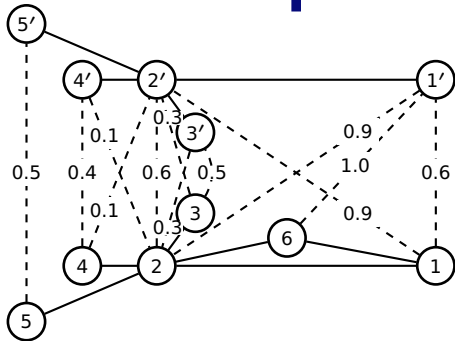


## Problem

maximize  $\sum_{i: e_i \in \mathcal{L}} x_i w_i + \sum_{i, j \in \mathcal{S}} x_i x_j$   
 subject to  $x$  is a matching

maximize  $\mathbf{w}^T \mathbf{x} + \frac{1}{2} \mathbf{x}^T \mathbf{S} \mathbf{x}$   
 subject to  $\mathbf{A} \mathbf{x} \leq \mathbf{e}$   
 $x_i \in \{0, 1\}$

# Example with overlap



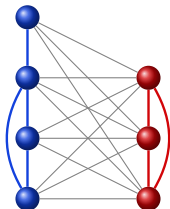
- $e_1=(2,2')$
- $e_2=(2,1')$
- $e_3=(2,3')$
- $e_4=(2,4')$
- $e_5=(1,2')$
- $e_6=(1,1')$
- $e_7=(3,2')$
- $e_8=(3,3')$
- $e_9=(4,2')$
- $e_{10}=(4,4')$
- $e_{11}=(5,5')$
- $e_{12}=(6,1')$

$$\mathbf{W} = \begin{pmatrix} 0.6 \\ 0.9 \\ 0.3 \\ 0.1 \\ 0.9 \\ 0.6 \\ 0.3 \\ 0.5 \\ 0.1 \\ 0.4 \\ 0.5 \\ 1.0 \end{pmatrix}$$

$$\mathbf{A} = \begin{pmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\mathbf{S} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

# An important special case



When  $\mathcal{L}$  contains everything possible (*the complete bipartite graph*), then  $\mathbf{S} = \mathbf{A} \otimes \mathbf{B}$  (*possibly permuted*).

**Important for discussion of related problems**

# Network alignment

NETWORK ALIGNMENT

maximize  $\alpha \mathbf{w}^T \mathbf{x} + \frac{\beta}{2} \mathbf{x}^T \mathbf{S} \mathbf{x}$

subject to  $\mathbf{A} \mathbf{x} \leq \mathbf{e}, x_i \in \{0, 1\}$

## History

- ▶ MAXIMUM COMMON SUBGRAPH
- ▶ PATTERN RECOGNITION
- ▶ ONTOLOGY MATCHING
- ▶ SIMILARITY FLOODING
- ▶ GRAPH SIMILARITY
- ▶ BIOINFORMATICS

## Sparse problems

Sparse  $\mathcal{L}$  largely ignored (exception Klau)

Our paper tackles that case explicitly

*Conte et al. Thirty years of graph matching, 2004.; Melnik et al. Similarity flooding, 2004; Blondel et al. SIREV 2004; Singh et al. RECOMB 2007; Klau, BMC Bioinformatics 10:S59, 2009.*

# Algorithms

1. SNOPT Relax, solve QP with SNOPT
2. ISORANK *(skipped)*
3. LP Convert to LP, relax, solve
4. TIGHTLP Improve the LP
5. BP *(skipped)*
6. MR *(skipped)*

**Large scale iterative methods skipped, see the paper.**

*Note The BP algorithm is based on max-product belief propagation for a large-scale approach. ISORANK uses a PageRank inspired heuristic, and MR uses a sequence of matching problems as a sub-gradient approach.*

Singh et al. RECOMB2007; Klau, 2009

# SNOPT

RELAXED NETWORK ALIGNMENT

$$\text{maximize } \alpha \mathbf{w}^T \mathbf{x} + \frac{\beta}{2} \mathbf{x}^T \mathbf{S} \mathbf{x}$$

$$\text{subject to } \mathbf{A} \mathbf{x} \leq \mathbf{e}, 0 \leq x_i \leq 1$$

- ▶ **S** is indefinite  
*(still NP-hard)*
- ▶ only local maximum,  
no upper-bound on  
the solution

**But  $\mathbf{x}$  is still fractional!** Use MAX-WEIGHT MATCHING with “weight” from the fractional  $\mathbf{x}$  to “round” the solution.

SNOPT	
Overlap	11,155
Weight	49,236 (100,886)

*Note* Took 6 days to compute, did not reach an optimum.



# LP

$$\mathbf{x}^T \mathbf{S} \mathbf{x} = \mathbf{e}^T \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & x_1 x_6 & 0 & x_1 x_8 & 0 & x_1 x_{10} & x_1 x_{11} & x_1 x_{12} \\ 0 & 0 & 0 & 0 & x_2 x_5 & 0 & x_2 x_7 & 0 & x_2 x_9 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & x_3 x_5 & 0 & x_3 x_7 & 0 & x_3 x_9 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & x_4 x_5 & 0 & x_4 x_7 & 0 & x_4 x_9 & 0 & 0 & 0 \\ 0 & x_5 x_2 & x_5 x_3 & x_5 x_4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & x_5 x_{12} \\ x_6 x_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & x_7 x_2 & x_7 x_3 & x_7 x_4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ x_8 x_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & x_9 x_2 & x_9 x_3 & x_9 x_4 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ x_{10} x_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ x_{11} x_1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ x_{12} x_1 & 0 & 0 & 0 & x_{12} x_5 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \mathbf{e}$$

$$\mathbf{y} \leq \mathbf{e}^T \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & y_{1,6} & 0 & y_{1,8} & 0 & y_{1,10} & y_{1,11} & y_{1,12} \\ 0 & 0 & 0 & 0 & y_{2,5} & 0 & y_{2,7} & 0 & y_{2,9} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & y_{3,5} & 0 & y_{3,7} & 0 & y_{3,9} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & y_{4,5} & 0 & y_{4,7} & 0 & y_{4,9} & 0 & 0 & 0 \\ 0 & y_{5,2} & y_{5,3} & y_{5,4} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & y_{5,12} \\ y_{6,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & y_{7,2} & y_{7,3} & y_{7,4} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ y_{8,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & y_{9,2} & y_{9,3} & y_{9,4} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ y_{10,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ y_{11,1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ y_{12,1} & 0 & 0 & 0 & y_{12,5} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \mathbf{e}; \quad \begin{matrix} y_{i,j} \leq x_i \\ y_{i,j} \leq x_j \end{matrix}$$

# LP

## NETWORK ALIGNMENT LP

$$\underset{\mathbf{x}, \mathbf{y}}{\text{maximize}} \quad \alpha \mathbf{w}^T \mathbf{x} + \frac{\beta}{2} \mathbf{e}^T \mathbf{Y}_S \mathbf{e}$$

$$\text{subject to} \quad \mathbf{A} \mathbf{x} \leq \mathbf{e}, x_i \in \{0, 1\}$$

$$\underbrace{\mathbf{Y}_S \leq \mathbf{X} \mathbf{S}}_{y_{i,j} \leq x_i}, \underbrace{\mathbf{Y}_S \leq \mathbf{S} \mathbf{X}}_{y_{i,j} \leq x_j}$$

- ▶ solve relaxed version
- ▶ more variables
- ▶ “big” LP solver required
- ▶ **upper bound** on solution

$\mathbf{x}$  still fractional, use  $\alpha \mathbf{w} + \beta/2 \mathbf{Y}_S \mathbf{x}$  as the “weight” for MAX-WEIGHT MATCHING rounding.

NETALIGNLP

Overlap 14,253

Weight 46,327 (96,297)

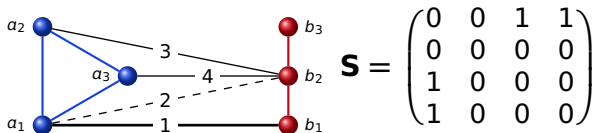
Upper bound on overlap 28660

*Note Took 1 day to compute*

*Klau, BMC Bioinformatics 10:S59, 2009.*

# TIGHTLP

Trivial upper bound on overlap  $\mathbf{e}^T \mathbf{S} \mathbf{e} / 2 = 892,655$



$\mathbf{x}$  must be a matching eventually.

$$\mathbf{x}^T \mathbf{S} \mathbf{x} = \sum x_i \mathbf{s}_i^T \mathbf{x} \leq \sum \mathbf{s}_i^T \mathbf{x} \leq \sum \max_{\text{match}}(\mathbf{s}_i, \mathbf{A}) = 344234.5$$

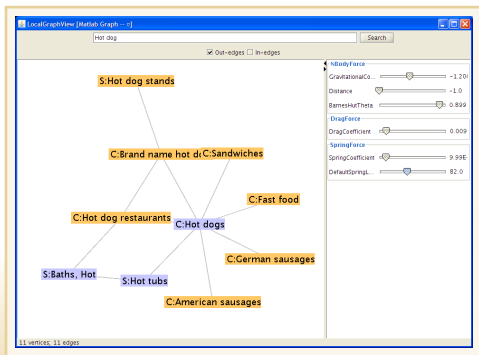
NETALIGN TIGHTLP
Overlap 17,251
Weight 46,270 (96,141)

Upper bound on overlap 17,571

Note Took 1 day to compute

# Matching results

LCSH	WC
Science fiction television series	Science fiction television programs
Turing test	Turing test
Maching learning	Machine learning
Hot tubs	Hot dog



# ***Graph similarity***

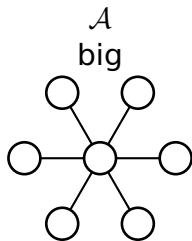
*Slide 21 of 29*

Network alignment

Graph similarity

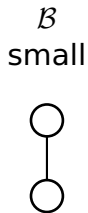
DAVID'S GENERAL WARNING: SPECULATIVE AND UNVERIFIED. SOME MATERIAL MAY CONFUSE ME

# Graph similarity



$\mathcal{L}$   
complete

"all edges"



**Make it a many to 1 matching**

# Graph similarity

## Problem

$$\begin{aligned} & \text{maximize}_{X_{r,u}} && \sum_{r,u} X_{r,u} W_{r,u} + \sum_{(r,u),(s,v) \in \mathcal{S}} X_{r,u} X_{s,v} \\ & \text{subject to} && \underbrace{\sum_u X_{r,u} \leq 1}_{\text{matching for } \mathcal{A} \rightarrow \mathcal{B}} \quad \forall r \end{aligned}$$

↔

$$\begin{aligned} & \text{maximize}_{\mathbf{x}} && \alpha \mathbf{w}^T \mathbf{x} + \frac{\beta}{2} \mathbf{x}^T \mathbf{S} \mathbf{x} \\ & \text{subject to} && \mathbf{A} \mathbf{x} = \mathbf{e} \end{aligned}$$

## Integer? Rounding?

*Note* For integer solutions, Yinyu Ye has a rounding procedure that would give you an approximation algorithm for the problem.



# HITS as graph similarity

## Idea

Given a graph  $\mathcal{G}$ , look at the similarity of  $\mathcal{G}$  to  $\textcircled{A} \rightarrow \textcircled{B}$

## Formulation

Largest eigenvalue of  $\mathbf{S} = \mathbf{A}^T \otimes \mathbf{B}^T + \mathbf{A} \otimes \mathbf{B}$  gives HITS scores,

maximize  $\mathbf{x}^T \mathbf{S} \mathbf{x}$

subject to  $\|\mathbf{x}\|_2 = 1$

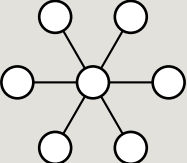



## Weaknesses

- ▶ if  $\mathbf{A}$  or  $\mathbf{B}$  is symmetric, then  $\mathbf{X}$  has rank 1
- ▶ no way to handle “sparse” case

**We can do the same thing too!**

*Should I have cited Kleinberg too? :- ) Blondel et al. SIREV 2004.*

# Examples

$A$	$B$	$X$
		$\begin{pmatrix} 0 & 1 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \\ 1 & 0 \end{pmatrix}$
		$\begin{pmatrix} 0 & 1 \\ 1 & 0 \\ 0 & 1 \\ 1 & 0 \\ 0 & 1 \end{pmatrix}$
$(\text{rand}(10) > 0.5)$	$(\text{rand}(10) > 0.5)$	$0.1\mathbf{e}\mathbf{e}^T$


Note Studied already? Close to **PATTERN RECOGNITION** formulation.  
 When  $B = \text{---}$  seems like **MAXCUT**, or finding a bipartition.

# Sports ranking

**Network alignment works on directed graphs too.**  
*(Just define squares carefully.)*

1. Align team networks against the dominance graph

$$\begin{pmatrix} 0 & 1 & \dots & 1 \\ 0 & 0 & \ddots & \vdots \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

2. Evaluate similarity against the dominance graph (allows ties)
3. Evaluate similarity against 

# Example

	Dav	UT	Cit	Cof	Wof	Sam	App	Elo	W.	Fur	UNC	GA
Dav		14	15	4	20	12	17	14	24	28	21	32
UT			6	11	7	9	4	12	7	27	14	13
Cit	18			12		25	4	2	12	13	7	15
Cof	5	9			2	7	6	8	15	9	15	13
Wof		3	5			11	6	3	15	14	12	7
Sam			9				11	13	12	8	14	23
App		6			5			19	6	4	1	12
Elo		1	2	6	7				6	8		9
W.		9		2	12	7	5	7		10	9	14
Fur		2				6					5	10
UNC							2	13				8
GA		4			21	7				16	9	

NETALIGN TIGHTLP

Overlap 64

Upper/Lower 739/197

Upper bound on overlap 65

*Note Thanks to Kathryn Pendings for sending the data on short notice!*

# Summary

## Network alignment

- ▶ Motivation from LCSH to Wikipedia
- ▶ QP, LP, and TightLP formulations
- ▶ Results (and problems) with LCSH to Wikipedia

## Graph similarity

- ▶ “Network alignment” inspired formulation
- ▶ A few small examples

## Sports ranking

- ▶ Interesting ideas for ranking teams?

*When you have a hammer, everything looks like a nail.*

